Linux Capabilities and Namespaces

Capabilities

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February 2023

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Rationale for capabilities

- Traditional UNIX privilege model divides users into two groups:
 - Normal users, subject to privilege checking based on UID and GIDs
 - Effective UID 0 (superuser) bypasses many of those checks
- Coarse granularity is a problem:
 - E.g., to give a process power to change system time, we must also give it power to bypass file permission checks
 - $\bullet \ \Rightarrow$ No limit on possible damage if program is compromised

[TLPI §39.1]

Rationale for capabilities

- Capabilities divide power of superuser into small pieces
 - 41 capabilities, as at Linux 6.2
 - Traditional superuser == process that has full set of capabilities
- Goal: replace set-UID-root programs with programs that have capabilities
 - Set-UID-root binary compromised ⇒ very dangerous
 - Compromise in binary with file capabilities ⇒ less dangerous
- Capabilities are not specified by POSIX
 - A 1990s standardization effort was ultimately abandoned
 - Some other implementations have something similar
 - E.g., Solaris, FreeBSD

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A selection of Linux capabilities

Capability	Permits process to
CAP_CHOWN	Make arbitrary changes to file UIDs and GIDs
CAP_DAC_OVERRIDE	Bypass file RWX permission checks
CAP_DAC_READ_SEARCH	Bypass file R and directory X permission checks
CAP_IPC_LOCK	Lock memory
CAP_FOWNER	<pre>chmod(), utime(), set ACLs on arbitrary files</pre>
CAP_KILL	Send signals to arbitrary processes
CAP_NET_ADMIN	Various network-related operations
CAP_SETFCAP	Set file capabilities
CAP_SETGID	Make arbitrary changes to process's (own) GIDs
CAP_SETPCAP	Make changes to process's (own) capabilities
CAP_SETUID	Make arbitrary changes to process's (own) UIDs
CAP_SYS_ADMIN	Perform a wide range of system admin tasks
CAP_SYS_BOOT	Reboot the system
CAP_SYS_NICE	Change process priority and scheduling policy
CAP_SYS_MODULE	Load and unload kernel modules
CAP_SYS_RESOURCE	Raise process resource limits, override some limits
CAP_SYS_TIME	Modify the system clock

More details: *capabilities*(7) man page and TLPI §39.2

Supporting capabilities

- To support implementation of capabilities, the kernel must:
 - Check process capabilities for each privileged operation
 - Cf. traditional check: is process's effective UID 0?
 - Provide system calls allowing a process to modify its capabilities
 - So process can raise (add) and lower (remove) capabilities
 - (Capabilities analog of set*id() calls)
 - Support attaching capabilities to executable files
 - When file is executed, process gains attached capabilities
 - (Capabilities analog of set-UID-root program)
- Implemented as follows:
 - Support for first two pieces available since Linux 2.2 (1999)
 - Support for file capabilities added in Linux 2.6.24 (2008)
 - (Delay due to design concerns rather than technical reasons)

[TLPI §39.4]

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Process and file capabilities

- Processes and (binary) files can each have capabilities
- Process capabilities define power of process to do privileged operations
 - Traditional superuser == process that has **all** capabilities
- File capabilities are a mechanism to give a process capabilities when it execs the file
 - Stored in security.capability extended attribute
 - (File metadata)

[TLPI §39.3]

Process and file capability sets

- Capability set: bit mask representing a group of capabilities
- Each **process**[†] has 3[‡] capability sets:
 - Permitted
 - Effective
 - Inheritable

†In truth, capabilities are a per-thread attribute ‡In truth, there are more capability sets

- An executable file may have 3 associated capability sets:
 - Permitted
 - Effective
 - Inheritable
- 1 Inheritable capabilities are little used; can mostly ignore

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Viewing process capabilities

• /proc/PID/status fields (hexadecimal bit masks):

- See <sys/capability.h> for capability bit numbers
 - Here: CAP_KILL (bit 5), CAP_SYS_ADMIN (bit 21)
- getpcaps(1) (part of libcap package):

```
$ getpcaps 4091
Capabilities for `4091': = cap_kill,cap_sys_admin+p
```

- More readable notation, but a little tricky to interpret
- Here, single '=' means inheritable + effective sets are empty
- capsh(1) can be used to decode hex masks:

```
$ capsh --decode=200020
0x000000000200020=cap_kill,cap_sys_admin
```

Modifying proces	s capabilities		
 Raising Syn Lowerin Syn Mostly, v 	onyms: drop , clea we'll defer discuss rious rules about	ing it to set) e moving it from set)	_
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Process permitted and effective capabilities

- Permitted: capabilities that process may employ
 - "Upper bound" on effective capability set
 - Once dropped from permitted set, a capability can't be reacquired
 - (But see discussion of execve() later)
 - Can't drop while capability is also in effective set
- Effective: capabilities that are currently in effect for process
 - I.e., capabilities that are examined when checking if a process can perform a privileged operation
 - Capabilities can be dropped from effective set and reacquired
 - Operate with least privilege....
 - Reacquisition possible only if capability is in permitted set

[TLPI §39.3.3]

File permitted and effective capabilities

- Permitted: a set of capabilities that may be added to process's permitted set during exec()
- Effective:
 <u>∧</u> a single bit that determines state of process's new effective set after exec():
 - If set, all capabilities in process's new permitted set are also enabled in effective set
 - Useful for so-called capabilities-dumb applications
 - If not set, process's new effective set is empty
- File capabilities allow implementation of capabilities analog of set-UID-root program
 - Notable difference: setting effective bit off allows a program to start in unprivileged state
 - Set-UID/set-GID programs always start in **privileged** state

[TLPI §39.3.4]

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3-3 3-8 -14 -18
-14
-18
-0
-30
-34
-39
-42
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Setting and viewing file capabilities from the shell

- setcap(8) sets capabilities on files
 - Requires privilege (CAP_SETFCAP)
 - E.g., to set CAP_SYS_TIME as a permitted and effective capability on an executable file:

```
$ cp /bin/date mydate
$ sudo setcap "cap_sys_time=pe" mydate
```

• getcap(8) displays capabilities associated with a file

```
$ getcap mydate
mydate = cap_sys_time+ep
```

- To list all files on the system that have capabilities, use: sudo filecap -a
 - filecap is part of the libcap-ng-utils package

[TLPI §39.3.6]

cap/demo_file_caps.c

- Display process capabilities
- Report result of opening file named in argv[1] (if present)

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cap/demo_file_caps.c

- All steps in demos are done from unprivileged user ID 1000
- Binary has no capabilities ⇒ process gains no capabilities
- open() of /etc/shadow fails
 - Because /etc/shadow is readable only by privileged process
 - Process needs CAP_DAC_READ_SEARCH capability

cap/demo_file_caps.c

```
$ sudo setcap cap_dac_read_search=p demo_file_caps
$ ./demo_file_caps /etc/shadow
Capabilities: = cap_dac_read_search+p
Open failed: Permission denied
```

- Binary confers permitted capability to process, but capability is not effective
- Process gains capability in permitted set
- open() of /etc/shadow fails
 - Because CAP_DAC_READ_SEARCH is not in effective set

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cap/demo_file_caps.c

```
$ sudo setcap cap_dac_read_search=pe demo_file_caps
$ ./demo_file_caps /etc/shadow
Capabilities: = cap_dac_read_search+ep
Successfully opened /etc/shadow
```

- Binary confers permitted capability and has effective bit on
- Process gains capability in permitted and effective sets
- open() of /etc/shadow succeeds

Notes for online practical sessions

- Small groups in breakout rooms
 - Write a note into Slack if you have a preferred group
- We will go faster, if groups collaborate on solving the exercise(s)
 - You can share a screen in your room
- I will circulate regularly between rooms to answer questions
- Zoom has an "Ask for help" button...
- Keep an eye on the #general Slack channel
 - Perhaps with further info about exercise;
 - Or a note that the exercise merges into a break
- When your room has finished, write a message in the Slack channel: "***** Room X has finished *****"
 - Then I have an idea of how many people have finished

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Shared screen etiquette

- It may help your colleagues if you use a larger than normal font!
 - In many environments (e.g., xterm, VS Code), we can adjust the font size with Control+Shift+"+" and Control+"-"
 - Or (e.g., emacs) hold down Control key and use mouse wheel
- Long shell prompts make reading your shell session difficult
 - Use PS1='\$ ' or PS1='# '
- Low contrast color themes are difficult to read; change this if you can
- Turn on line numbering in your editor
 - In vim use: :set number
 - In emacs use: M-x display-line-numbers-mode <RETURN>
 M-x means Left-Alt+x
- For collaborative editing, relative line-numbering is evil....
 - In vim use: :set nornu
 - In emacs, the following should suffice:

```
M-: (display-line-numbers-mode) <RETURN>
M-: (setq display-line-numbers 'absolute) <RETURN>
```

M-: means Left-Alt+Shift+:

Using tmate in in-person practical sessions

In order to share an X-term session with others, do the following:

• Enter the command *tmate* in an X-term, and you will see the following:

```
$ tmate
...
Connecting to ssh.tmate.io...
Note: clear your terminal before sharing readonly access
web session read only: ...
ssh session read only: ...
web session: ...
ssh session: ssh SOmErAnDOm5Tr1Ng@lon1.tmate.io
```

- Share the last "ssh" string with your colleagues via Slack or another channel
- Your colleagues should paste that string into an X-term...
 - After that, you will be sharing an X-term session in which anyone can type

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Exercises

Compile and run the cap/demo_file_caps program, without adding any capabilities to the file, and verify that the process has no capabilities when it executes the binary:

```
$ cc -o demo_file_caps demo_file_caps.c -lcap
```

- The string "=" means, all capability sets empty.)
- 2 Now make the binary set-UID-root:

```
$ sudo chown root demo_file_caps # Change owner to root
$ sudo chmod u+s demo_file_caps # Turn on set-UID bit
$ ls -l demo_file_caps # Verify
-rwsr-xr-x. 1 root mtk 8624 Oct 1 13:19 demo_file_caps
```

Question Run the binary and verify that the process gains all capabilities. (The string "=ep" means "all capabilities in the permitted and effective sets".)

[Exercise continues on next slide]

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- Take the existing set-UID-root binary, add a permitted capability to it, and set the effective capability bit:
 - \$ sudo setcap cap_dac_read_search=pe demo_file_caps
- When you now run the binary, what capabilities does the process have?
- Suppose you assign empty capability sets to the binary. When you execute the binary, what capabilities does the process then have?
 - \$ sudo setcap = demo_file_caps
- Use the setcap -r command to remove capabilities from the binary and verify that when run, it once more grants all capabilities to the process.

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Capabilities-dumb and capabilities-aware applications

- Capabilities-dumb application:
 - (Typically) an existing set-UID-root binary whose code we can't change
 - Thus, binary does not know to use capabilities APIs (Binary simply uses traditional set*uid() APIs)
 - But want to make legacy binary less dangerous than set-UID-root
- Converse is capabilities-aware application
 - Program that was written/modified to use capabilities APIs
 - Set binary up with file effective capability bit off
 - Program "knows" it must use capabilities APIs to enable effective capabilities

Adding capabilities to a capabilities-dumb application

To convert existing set-UID-*root* binary to use file capabilities:

- Setup:
 - Binary remains set-UID-root
 - Enable a subset of file permitted capabilities + set effective bit on
 - (Note: code of binary isn't changed)
- Operation:
 - When binary is executed, process gets (just the) specified subset of capabilities in its permitted and effective sets
 - IOW: file-capabilities override effect of set-UID-*root* bit, which would normally confer **all** capabilities to process
 - Process UID changes between zero and nonzero automatically raise/lower process's capabilities
 - (Covered in more detail later)

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Textual representation of capabilities

- Both setcap(8) and getcap(8) work with textual representations of capabilities
 - Syntax described in cap_from_text(3) man page
- Strings read left to right, containing space-separated clauses
 - (The capability sets are initially considered to be empty)
- Clause: caps-list operator flags [operator flags] ...
 - caps-list is comma-separated list of capability names, or all
 - *operator* is =, +, or -
 - flags is zero or more of p (permitted), e (effective), or i (inheritable)

Textual representation of capabilities

- Operators:
 - = operator:
 - Raise capabilities in sets specified by flags; lower those capabilities in remaining sets
 - caps-list can be omitted; defaults to all
 - flags can be omitted ⇒ clear capabilities from all sets
 ⇒ Thus: "=" means clear all capabilities in all sets
 - + operator: raise capabilities in sets specified by flags
 - operator: lower capabilities in sets specified by flags
- Clause can contain multiple [operator flags] parts:
 - E.g., cap_sys_time+p-i
- What does "=p cap_kill,cap_sys_admin+e" mean?
 - All capabilities in permitted set, plus CAP_KILL and CAP_SYS_ADMIN in effective set

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Exercises

- What capability bits are enabled by each of the following text-form capability specifications?
 - "=p"
 - "="
 - "cap_setuid=p cap_sys_time+pie"
 - "cap_kill=p = cap_sys_admin+pe"
 - "cap_chown=i cap_kill=pe cap_setfcap,cap_chown=p"
 - "=p cap_kill-p"
- 2 The program cap/cap_text.c takes a single command-line argument, which is a text-form capability string. It converts that string to an in-memory representation and then iterates through the set of all capabilities, printing out the state of each capability within the permitted, effective, and inheritable sets. It thus provides a method of verifying your interpretation of text-form capability strings. Try supplying each of the above strings as an argument to the program (remember to enclose the entire string in quotes!) and check the results against your answers to the previous exercise.

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Transformation of process capabilities during exec

During execve(), process's capabilities are transformed:

```
P'(perm) = F(perm) & P(bset)

P'(eff) = F(eff) ? P'(perm) : 0
```

- P() / P'(): process capability set before/after exec
- F(): file capability set (of file that is being execed)
- New permitted set for process comes from file permitted set ANDed with capability bounding set (bset)
 - \triangle Note that P(perm) has no effect on P'(perm)
- New effective set is either 0 or same as new permitted set
- Transformation rules above are a simplification that ignores process+file inheritable sets and process ambient set
 - In most cases, those sets are empty (i.e., 0)

Transformation of process capabilities during exec

- Commonly, process bounding set contains all capabilities
- Therefore transformation rule for process permitted set:

```
P'(perm) = F(perm) & P(bset)
```

commonly simplifies to:

$$P'(perm) = F(perm)$$

[TLPI §39.5]

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The capability bounding set

- Per-process attribute (actually: per-thread)
- A "safety catch" to limit capabilities that can be gained during exec
 - Limits capabilities that can be granted by file permitted set
 - Limits capabilities that can be added to process inheritable set (later)
- Use case: remove some capabilities from bounding set to ensure process never regains them on execve()
 - E.g., *systemd* reduces bounding set before executing some daemons
 - Guarantees that daemon can never get certain capabilities
 - cap/reduced_bounding_set_procs.sh displays list of processes that have a reduced bounding set

The capability bounding set

- Inherited by child of fork(), preserved across execve()
 - init starts with capability bounding set containing all capabilities
- Two methods of getting:
 - prctl() PR_CAPBSET_READ (for self)
 - Higher-level *libcap* API: cap_get_bound(3)
 - /proc/PID/status CapBnd entry (any process)
- Can (irreversibly) drop capabilities from bounding set
 - prctl() PR_CAPBSET_DROP
 - Requires CAP_SETPCAP effective capability
 - Doesn't change permitted, effective, and inheritable sets
 - Higher-level libcap API: cap_drop_bound(3)

[TLPI §39.5.1]

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